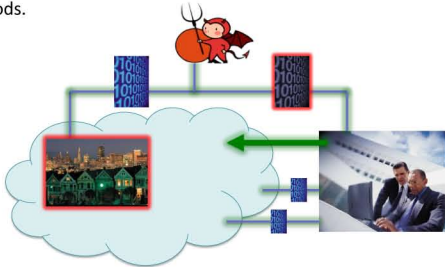


## Goals

- Reliable detection of bad-data injection attacks that are potentially undetectable by conventional methods, by using perturbation-based methods.



- Accurate modeling of effects on the power system of perturbations implemented to detect attacks.
- Improved understanding of full taxonomy of attacks that now are potentially undetectable by conventional methods.

## Fundamental Questions/Challenges

- Is there an iterative algorithm, based on application of perturbations to an electrical system and observation of the results, for identifying “bad-data” attacks without trusting any of the received measurements as valid?
- Does a cyber-physical system have inherent properties that can be used to undermine a cyber-only attack on it?
- What advantages against an attacker who wishes to remain undetected can a defender gain from physical assets and full system knowledge?

## Research Results

- We have presented and demonstrated an algorithm that produces attack vectors that simultaneously target arbitrary, specific measurements *and* result in altered measurements that satisfy the full AC power flow model equations—rendering the attack undetectable by conventional methods.

TRUE DATA VALUES

ATTACK POWER FLOW VALUES

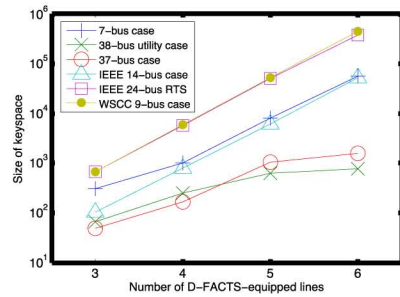
| Bus Num. | V(pu) | Theta(deg) | Pinj(pu) | Qinj(pu) | Bus Num. | V(pu) | Theta(deg) | Pinj(pu) | Qinj(pu) |
|----------|-------|------------|----------|----------|----------|-------|------------|----------|----------|
| 1        | 1.04  | 0          | 0.717    | 0.270    | 1        | 1.04  | 0          | 0.716    | 0.270    |
| 2        | 1.025 | 9.279      | 1.63     | 0.066    | 2        | 1.025 | 9.280      | 1.63     | 0.066    |
| 3        | 1.025 | -4.664     | 0.85     | -0.109   | 3        | 1.025 | 4.664      | 0.850    | -0.109   |
| 4        | 1.026 | -2.217     | 0.0      | 0.0      | 4        | 1.026 | -2.217     | -0.001   | -0.640   |
| 5        | 0.996 | -3.989     | -1.25    | -0.5     | 5        | 0.996 | -3.989     | -1.25    | -0.500   |
| 6        | 1.013 | -3.688     | -0.9     | -0.3     | 6        | 1.070 | -4.173     | -0.9     | 0.7      |
| 7        | 1.026 | 3.718      | 0.0      | 0.0      | 7        | 1.026 | 3.720      | 0.000    | 0        |
| 8        | 1.016 | 0.726      | -1.0     | -0.35    | 8        | 1.016 | 0.727      | -1       | -0.350   |
| 9        | 1.032 | 1.966      | 0.0      | 0.0      | 9        | 1.032 | 1.966      | 0.010    | -0.345   |

(Values given by attack algorithm example on the WSCC 9-bus test system)

- The attack algorithm is essentially the reverse of the typical iterative power flow solution:
  - Rather than set one bus as the slack and solve for  $V$  &  $\Theta$  (PQ buses) or  $Q$  &  $\Theta$  (PV buses) at all other points, we chose all buses over which the attacker does not have control as effectively slack buses—fixing their  $V$  &  $\Theta$  and solving for  $P$  &  $Q$ .
  - What that effectively does is allow the adversary to find a physically viable attack vector that will change only specifically targeted measurements—such as sending of bad data showing the voltage at a particular bus to be too high (demonstrated in [1])—plus some measurements at adjacent “boundary” buses.
  - Those boundary bus modifications are needed to maintain the physically required power balance, and are analogous to the requirement that the defender be able to protect a set of “basic measurements” in order to detect attacks under the DC model.

## Research Results (continued)

- Initial feasibility analysis of the perturbation-based approach indicates potential utility. Computational experiments were run using six test systems (ranging from 7 to 38 buses), with D-FACTS devices on select lines used as the perturbation mechanism.
  - Possible spaces of perturbations (or “keys”) were large enough that a randomly chosen key was effectively unguessable by the adversary.
  - The sizes of those possible spaces (or “keyspaces”) were still sufficiently large even with the constraint that the effected changes had to be larger than some noise level.



## Broader Impact

- Rather than use the perturbation-based method to detect attacks reactively based on suspicion when something is amiss, one could also use this method to check for attacks preventatively on some continual, routine basis.
- While the work focuses on the impact of D-FACTS devices, it could be extended to work for other physical perturbations, including those continuously ongoing in the power grid during normal operations.

## Publications

- K. R. Davis, K. L. Morrow, R. Bobba, E. Heine. “Power Flow Cyber Attacks and Perturbation-Based Defense.” To be presented at *SmartGridComm 2012*.
- K. L. Morrow, E. Heine, K. M. Rogers, R. B. Bobba, T. J. Overbye. “Topology Perturbation for Detecting Malicious Data Injection.” In *Proceedings of the 45<sup>th</sup> Hawaii International Conference on Systems Sciences (HICSS '12)*, Maui, Hawaii, January 2012.

## Interaction with Other Projects

- Presented demo at 2011 TCIPG Industry Workshop (Urbana, IL) using DETER experiment and RTDS simulation to demonstrate detection of a compromised set of power flow measurements.

## Future Efforts

- Continue work on characterization of possible attack space: how the adversary would implement a given attack, and what that attack would look like to the defender. Can attacks be uniquely identified?
- More work to be done on modeling the operational effects of the proposed perturbation-based defense.
- Design an algorithm that uses perturbation key sequences to quickly detect and accurately determine the specific data injection attack being performed.

