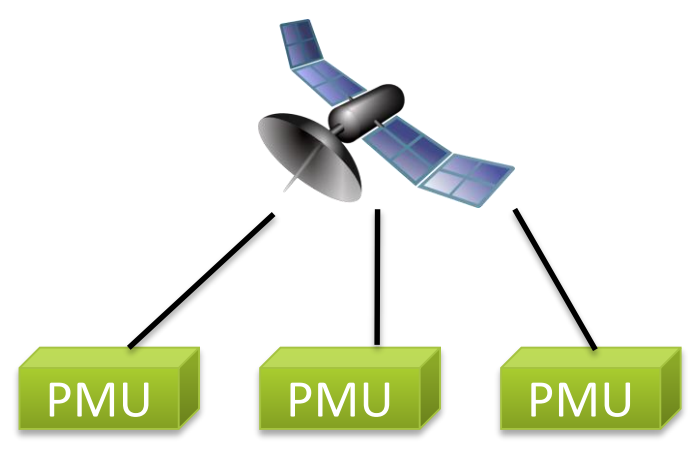
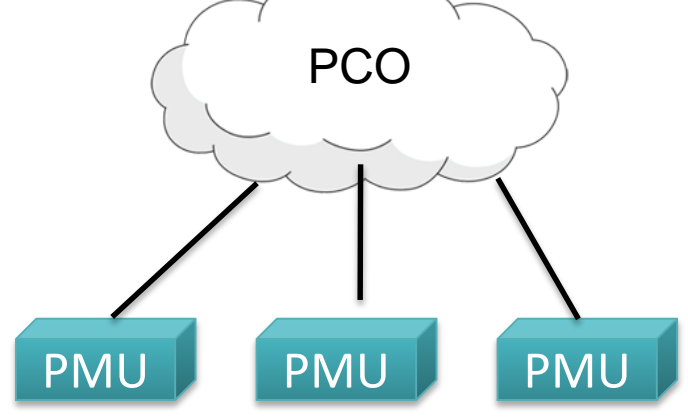


GOALS

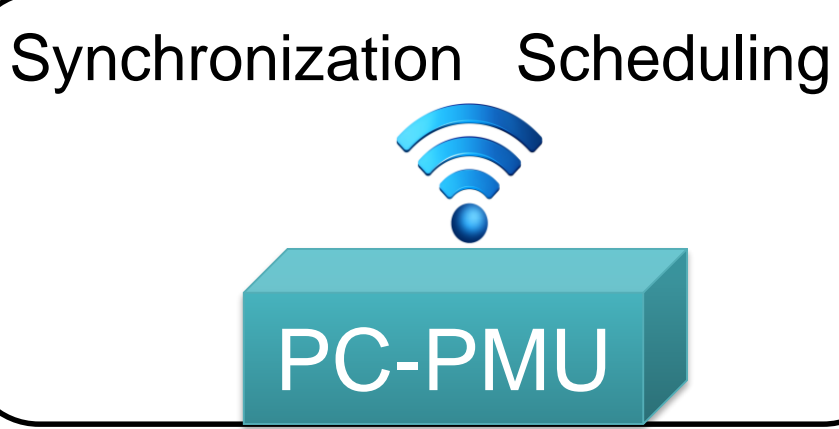


+ Scheduling

=



=



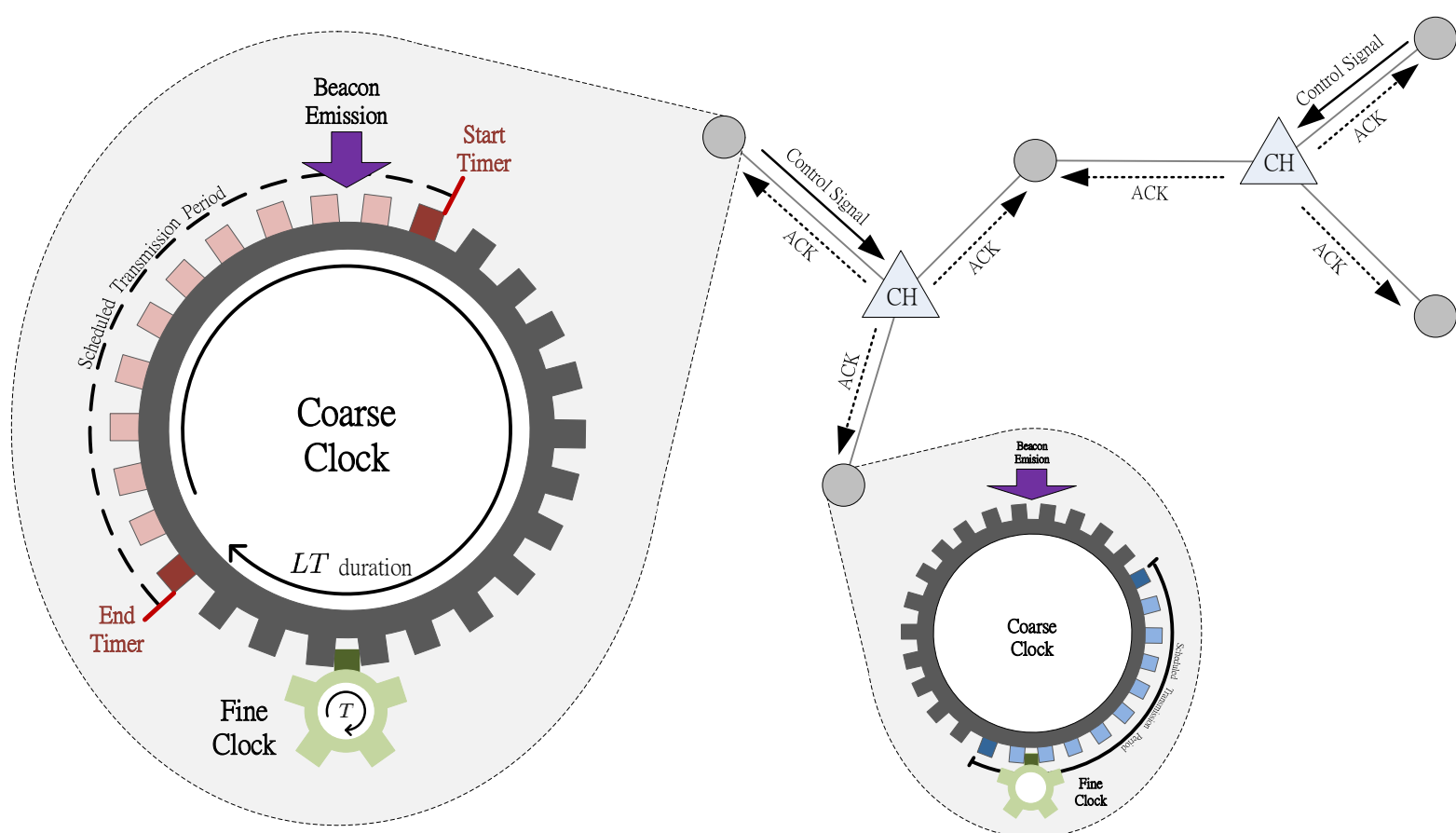
Accurate network timing is required for Phasor Measurement Units (PMU) or other sensors.

- GPS spoofing is a threat.
- PTP & NTP:
 - Not encrypted (man in the middle attack).
 - Centralized (not scalable; central point of failure).
- PCO-Synchronization + Scheduling:
 - Resilient to node failures.
 - Self-healing (inspired by biological networks).
 - Physical-layer secure signaling.
 - Decentralized design → scalable.
 - Accurate timing due to time of flight estimation.
 - Application: monitor and control for the distribution grid.

INTRODUCTION TO THE PULSESS PROTOCOL (PULSE-COUPLED SYNCHRONIZATION AND SCHEDULING)

Define start and end coarse clocks (length L) driven by the same fine clock.

$$\begin{aligned} \Phi_v(t) &= s_v(t) + \varphi_v(t) \pmod{L} \\ \Phi'_v(t) &= e_v(t) + \varphi_v(t) \pmod{L} \end{aligned}$$



- Each time a clock reaches L , it sends a pulse.
- Cluster heads (CH) in range spread the pulse locally by acknowledging it.
- The CHs adjust their own fine clocks, taking into account the estimated signal traveling time $\hat{t}_{v,c}$, i.e., when receiving the pulse at time $r_v^{(s)}$.

$$\hat{\varphi}_c = \varphi_c(r_v^{(s)}) - \frac{\hat{t}_{v,c}}{T}$$

$$\varphi_c(r_v^{(s)+}) = \begin{cases} \varphi_c(r_v^{(s)}), & \text{if } 1 - \hat{\varphi}_c \pmod{T} \leq \delta_{ref} \\ \min((1 + \alpha)\hat{\varphi}_c, 1) + \frac{\hat{t}_{v,c}}{T} \pmod{T}, & \text{else} \end{cases}$$

- The refractory period δ_{ref} is set based on the noise in time of arrival of a pulse:

$$\delta_{ref} \propto \sigma_{v,c}^2 = (SNR_{(v,c)} \cdot \bar{F}^2)^{-1}$$

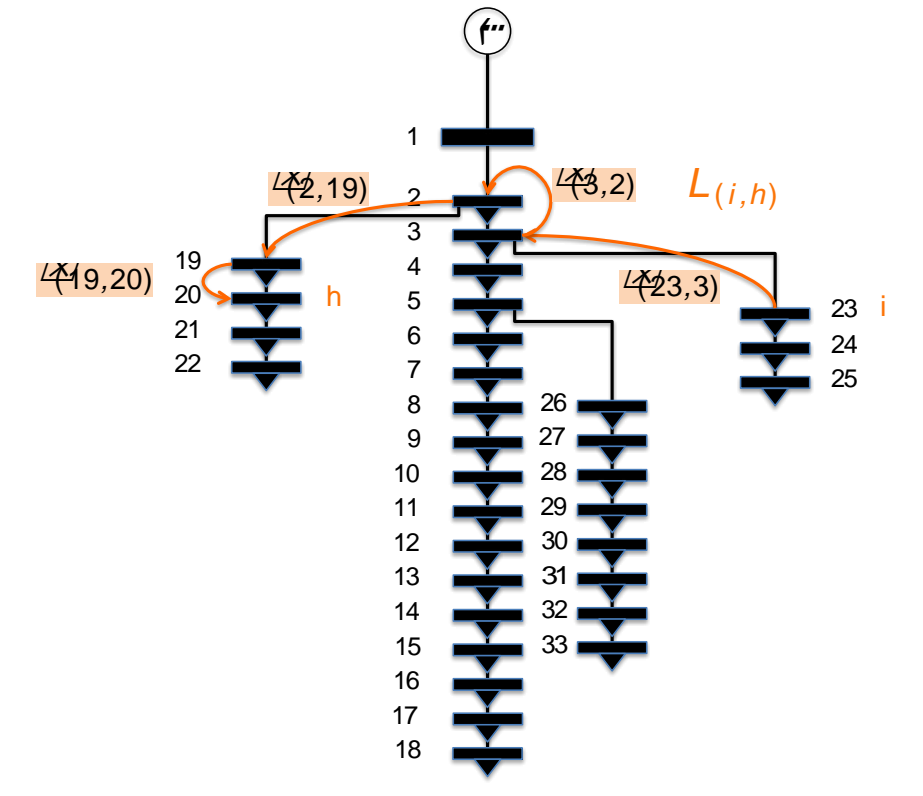
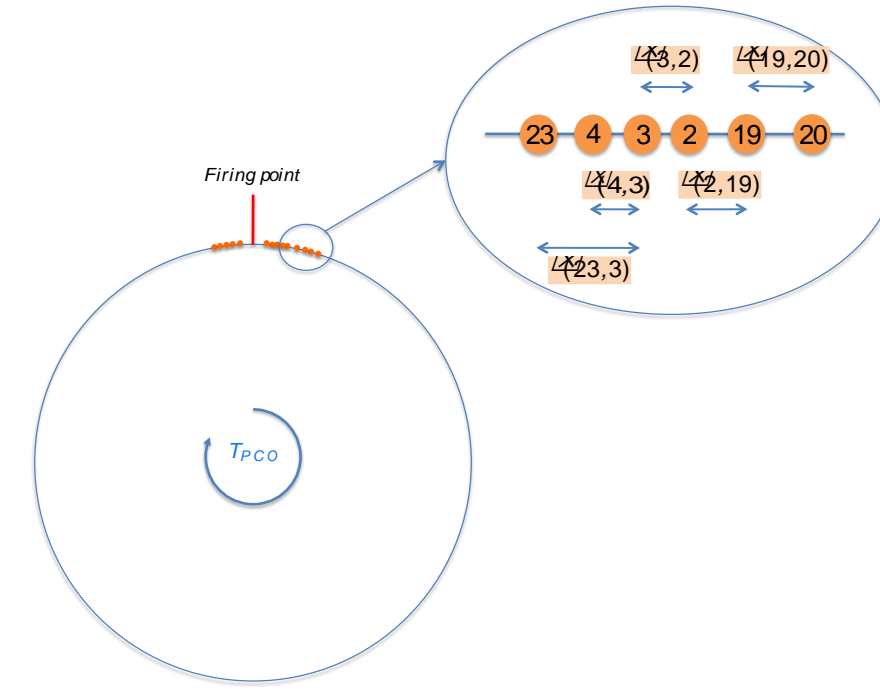
Thus, the accuracy depends on the signal-to-noise ratio of the link and mean squared bandwidth.

- Each **node** that receives an acknowledgment updates its fine clock analogously.
- In addition, when a node is receiving a start pulse from its successor, it updates its coarse clock.
- The update depends on the positions of the nodes' predecessors; the nodes' demand for communication, D_v ; and an intended guard space, δ .

$$\begin{aligned} s_v^{target}(t) &= \frac{(D_v + \delta)}{(D_v + 2\delta)} e_{pre(v)}(t) + \frac{\delta}{(D_v + 2\delta)} s_{suc(v)}(t) \\ e_v^{target}(t) &= \frac{(D_v + \delta)}{(D_v + 2\delta)} s_{suc(v)}(t) + \frac{\delta}{(D_v + 2\delta)} e_{pre(v)}(t) \end{aligned}$$

RESEARCH RESULTS

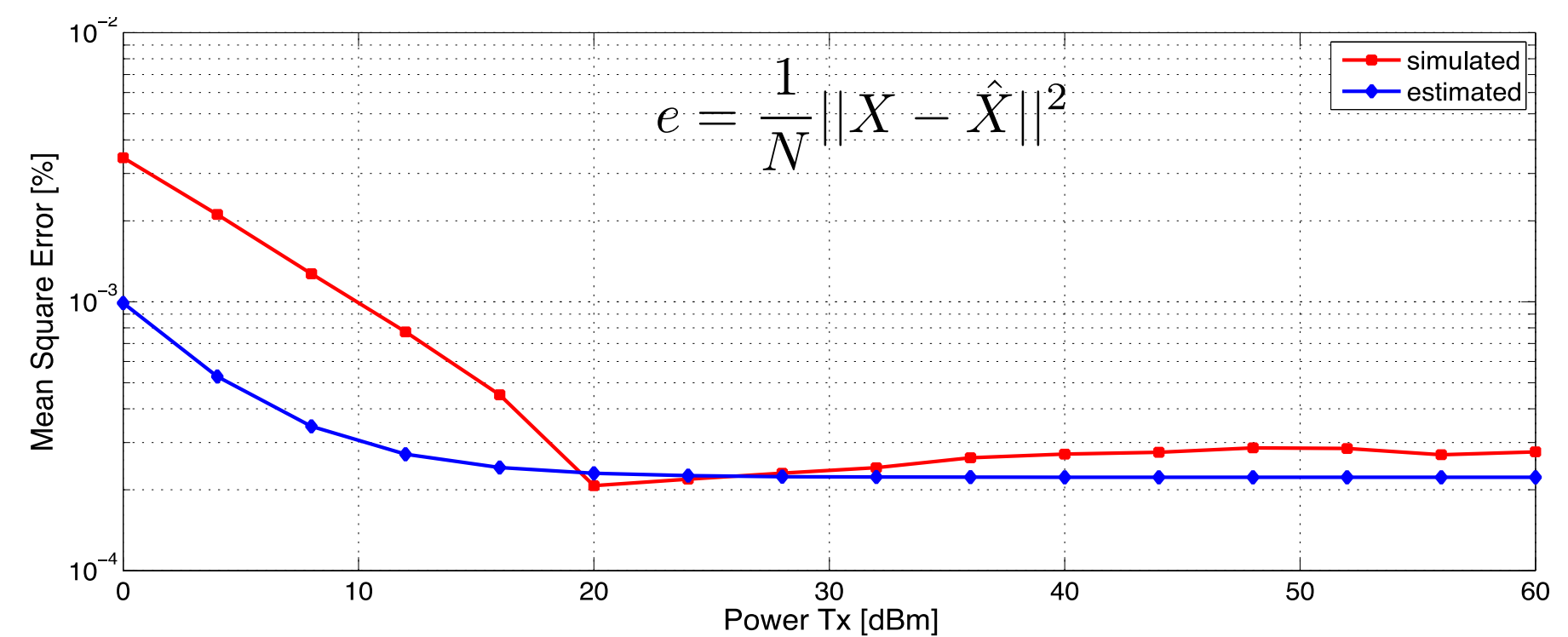
PCO Timing Error Analysis



Lemma (Convergence fixed point) $\Delta\Phi_{(i,h)}(t) \xrightarrow{t \rightarrow \infty} \sum_{(k,m) \in L(i,h)} \tau_{(k,m)}$

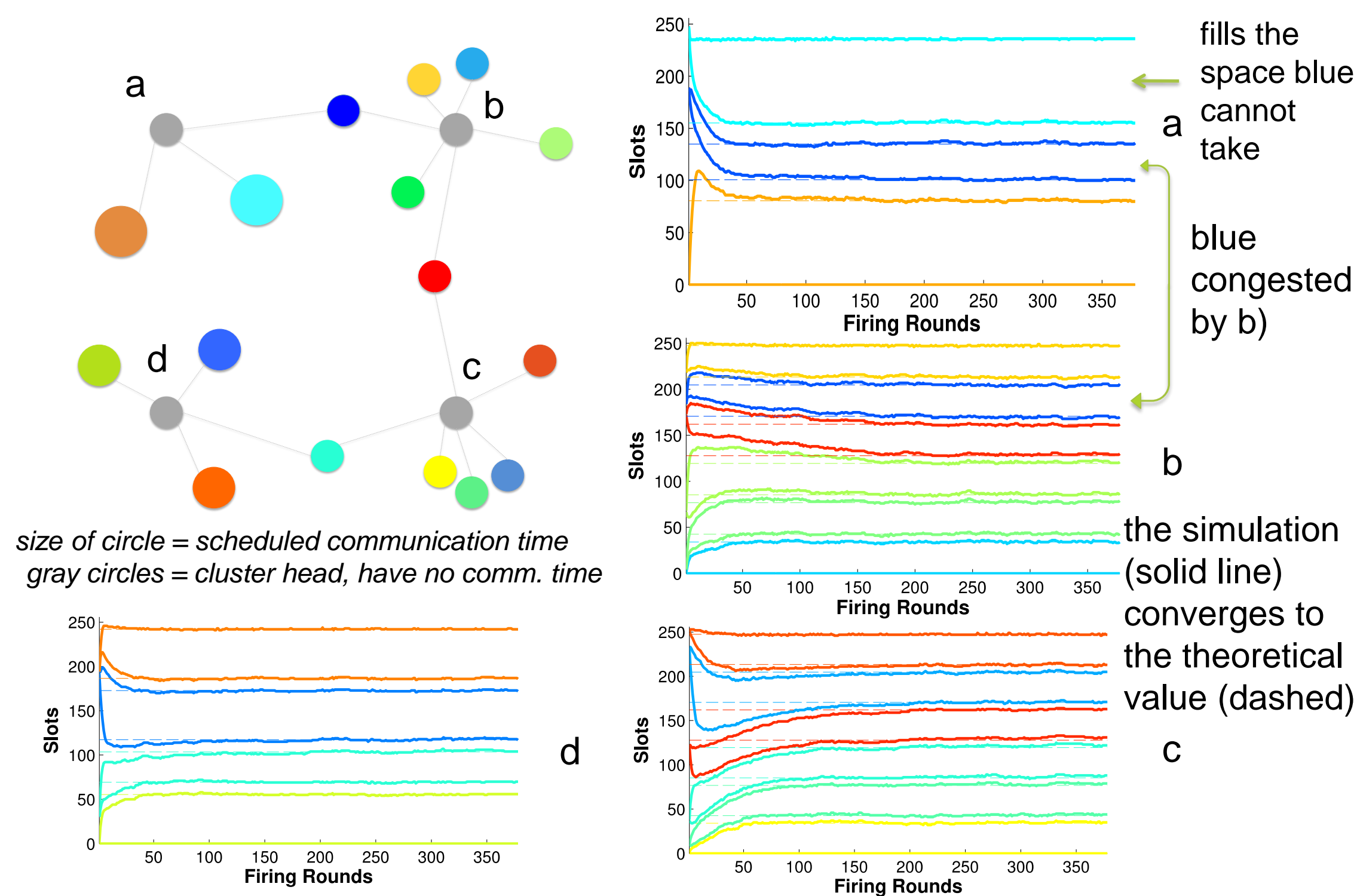
Simulation: IEEE 33 bus, 1 PC-PMU per branch. Power-line communication (band around 300kHz +/- 100kHz). Losses on the line 40dB/km, average distance 100 m, coupling factor $\alpha=0.04$, 170 iterations, transmission power 0-60dBm.

Mean squared error 0.023% @ 30dBm



Scheduling convergence of PulseSS for clustered networks

Simulation: 4 clusters; each node has the same demand; 250 slots; 370 iterations; transmission power 30dBm; other parameters as above.



BROADER IMPACT

- A decentralized self-healing radio protocol to support synchronization and scheduling can reduce vulnerabilities due to possible spoofing and jamming of GPS signals and other master/slave network synchronization protocols.
- It can make NICS scalable and easy to deploy because of reliable timing and scheduling of information flow.

FUTURE EFFORTS

- Complete the analysis of convergence speed.
- Study compatibility of PulseSS as a wake-up radio.
- Develop and test the PulseSS hardware implementation.
 - Test with microcontroller running TinyOS.
 - Test with FPGA, allowing direct access to the physical layer.